

## Loop (with Counter)

- We can loop for ~~specific~~ specific number of times using **DJNZ** instruction

**DJNZ Counter, TARGET**

- This instruction will <sup>register OR port</sup> Decrement the Counter then check if it ~~is~~ not equal Zero then Jump to target otherwise, ~~we~~ Continue normal flow of execution

Mov R2, ~~10~~ 10 } Loop for 10 times  
 Loop: DJNZ R2, Loop

- 8051 Registers are 8-bit So we can loop to 255 if we need more we should use nested loop

Example: Loop for 700 times

Mov R1, ~~7~~ 7  
 Loop1: Mov R2, ~~100~~ 100  
 Loop2: DJNZ R2, Loop2 → inner  
 DJNZ R1, Loop1 → outer

Conditional	Jump	Jump
JZ	target	→ if A = 0
JNZ	target	→ if A ≠ 0
DJNZ	R0, target	
CJNE	R0, <del>80</del> 80, target	→ if R0 ≠ 80
JC	target	→ if Carry = 1
JNC	target	→ if Carry = 0
JB	P0.1, target	→ if bit(P0.1) = 1
JNB	P1.0, target	→ if P1.0 = 0
JBC	C, target	→ if C = 1

## JBC Instruction

- This instruction check if the bit = 1
- If that is true Jump and Clear that bit

## CJNE instruction (slide 222)

- This instruction is used to compare two values
- The comparison is performed by subtracting the two values without change values in registers

- It only change the carry

**CJNE A, R0, Not-EQUAL**

NOT-EQUAL: ~~JNC~~ <sup>→ jump not carry</sup> A - R0

- If the Carry = 1 that mean [A - R0] → result is negative → A < R0
- If the Carry = 0 A > R0

## Un Conditional Jump

### Short Jump

**SJMP target**

- It takes only 2 Bytes
- the first byte for opcode
- the second byte is a relative

Address (-128 → 127)  
 Address    OP Code    Instruction  
 0007    XX 06    SJMP target  
 0009    XX 00    NOP  
 000F    target: NOP  
 0F - 09 = 06

- لاحظ أن يتغير حساب الـ relative address  
بالاعتماد على عنوان الأمر التالي  
المضافة إلى عنوان الذي يتغير إلى انتقال إليه

~~000F - 0009~~

$$000F - 0009 = 06$$

- لابد أن يكون الفرق في حدود 127, 128

- Short Jump نوع Conditional Jump

Address

0002

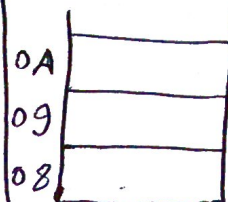
0005

Instruction

LCALL XYZ

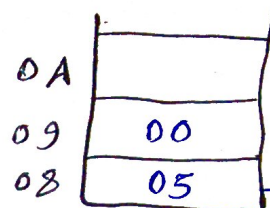
Nop

Stack



SP = 07

after LCALL



SP = 09

Low First

## 2 Long Jump

- It takes 3 bytes
- The first byte for opcode
- The other two bytes is the absolute Address of the target

LJMP target

Look At slide 117

- pop

- push

بتنظيف عتمة

## The CALL Instructions

### 1 Long Call (LCALL target)

- 3-bytes instruction
- there are two bytes for address

~~Short Call~~

### 2 Absolute Call (ACALL target)

- 2-bytes instruction
- 11 bits are used for address

- The CALL instruction Store the Content of Program Counter in Stack before Jump

- The RET instruction is used to POP [PC] from stack and back to instruction after call



## Delay in 8051

- Delay depends on

① Crystal frequency

② Design of 8051

For 8051

\* Crystal freq. = 11.0592 MHz

\* Clocks per Machine Cycle

- depend on Design

- Default = 12

To Compute machine Cycle

$$\text{machine freq.} = \frac{\text{Crystal Freq.}}{\text{Clocks per Machine Cycle}}$$

$$\text{machine cycle} = \frac{1}{\text{machine freq.}}$$

Common no. of machine cycles  
for Common Instruction

MOV	1
DEC	1
INC	1
DJNZ	2
LJMP	2
SJMP	2
NOP → no operation	1
MUL AB	4
PUSH	2
POP	2
RET	2

Look At Examples at  
Slides 123 → 127

Example

Write Delay Subroutine to introduce  
100 ms delay

$$\begin{aligned}\text{Machine freq.} &= \frac{11.0592}{12} \\ &= 921.6 \text{ KHz} \\ \text{machine cycle} &= 1/921.6 \text{ KHz} \\ &= 1.085 \mu\text{s}\end{aligned}$$

The required no. of machine cycles

$$\text{no. cycle} = \frac{100 \text{ ms}}{1.085 \mu\text{s}}$$

$$= 92165.89$$

$$\approx 92166 \text{ cycles}$$

We try to use two nested  
Loop

Delay: MOV R1, ~~200~~ → X

Outer: MOV R2, ~~228~~ → Y

inner: DJNZ R2, inner

DJNZ R1, outer

RET

Cycle  
1  
1  
1  
1  
1  
1

MOV → 1

MOV → 1 \* X

DJNZ → 1 \* X \* Y

DJNZ → 2 \* X

RET → 2

$$92166 = 3 + 3X + 2XY$$

Assume X = 150

$$Y \approx 305 > 255$$

Assume X = 200

$$Y \approx 229$$